## CS154: Homework #3

Due: Wednesday, July 25, 2012 by 5PM

**Problem 1.** Let L be a regular language over some alphabet  $\Sigma$ . Prove that the language T, defined as,

 $\mathbf{T} = \{ w \in \Sigma^* \mid \text{for some string } x, \, wx \in \mathbf{L} \}$ 

is also regular.

**Note:** Simply describing the DFA is *not* sufficient. You need to give a *formal* proof which has two sides, similar to that given in Lecture 1.

**Problem 2.** Let  $\mathbf{L} \subset \{0,1\}^*$  be the language of all strings such that there are two 0's separated by a number of positions that is a **non-zero** multiple of 5. Prove that any DFA for this language must have **at** least  $2^5$  states.

**Problem 3.** Consider two regular languages  $\mathbf{L}_1$  and  $\mathbf{L}_2$  defined over the *same* alphabet  $\Sigma$ . Let the *shuffle* of  $\mathbf{L}_1$  and  $\mathbf{L}_2$  be the language consisting of strings

 $\{w \mid w = a_1 b_1 a_2 b_2 \dots a_k b_k, \text{ where } a_1 \dots a_k \in \mathbf{L}_1 \text{ and } b_1 \dots b_k \in \mathbf{L}_2 \text{ and each } a_i, b_i \in \Sigma\}$ 

In other words, the shuffle consists of strings of equal length from  $\mathbf{L}_1$  and  $\mathbf{L}_2$  whose alphabets are **strictly alternated**. Prove that the shuffle of  $\mathbf{L}_1$  and  $\mathbf{L}_2$  is a regular language as well.

**Hint:** Let  $\Sigma' = \Sigma \times \Sigma$  and consider the homomorphism  $h : \Sigma' \to \Sigma$  defined as h(a, b) = a.

Problem 4. Consider the transition table of a DFA,

		0	1
÷	A	B	A
	B	A	C
	C	D	B
	$D^*$	D	A
	E	D	F
	F	G	E
	G	F	G
	H	G	D

- a) Draw the table of distinguishabilities for this automaton.
- b) Construct the minimum-state equivalent DFA.

Problem 5. Consider the CFG G defined by the productions

$$S \to aSbS \mid bSaS \mid \varepsilon$$

Prove that L(G) is the set of all strings with an equal numbers of a's and b's.