

# CS154: Homework #2

Due: Wednesday, July 18, 2012 by 5PM

**Problem 1.** For any DFA, we extend the transition function  $\delta$  by breaking the input string  $\mathbf{w} = \mathbf{xa}$  during the inductive step, where  $\mathbf{x}$  is any string followed by a single symbol  $\mathbf{a}$ . However, we informally think of  $\delta$  as describing what happens along a path with a certain string of symbols, and if so, then it should not matter how we break the input string. Show that in fact,

$$\delta(q, \mathbf{xy}) = \delta(\delta(q, \mathbf{x}), \mathbf{y})$$

for any state  $q$  and strings  $\mathbf{x}$  and  $\mathbf{y}$ .

**Problem 2.** Give DFA's accepting the following languages over the alphabet  $\{0, 1\}$ .

- The set of all strings such that each block of **four consecutive** symbols contains **at least two** 0's.
- The set of strings such that the number of 0's is divisible by 3, **and** the number of 1's is divisible by 3.

**Problem 3.** Consider the following  $\varepsilon$ -NFA.

	$\varepsilon$	$a$	$b$	$c$
$\rightarrow p$	$\{q, r\}$	$\emptyset$	$\{q\}$	$\{r\}$
$q$	$\emptyset$	$\{p\}$	$\{r\}$	$\{p, q\}$
$r^*$	$\emptyset$	$\emptyset$	$\emptyset$	$\emptyset$

Recall that starred states denote accept/final states.

- Compute the  $\varepsilon$ -closure of each state.
- Give all the strings of length three or less accepted by the automaton.
- Convert the automaton to a DFA.

**Problem 4.** Write regular expressions for the following languages. In all parts the alphabet is  $\{0, 1\}$ .

- $\{w \mid w \text{ starts with } 0 \text{ and has odd length, or starts with } 1 \text{ and has even length}\}$ .
- $\{w \mid w \text{ has at most one pair of consecutive } 1\text{'s}\}$ .
- $\{w \mid \text{the number of } 0\text{'s in } w \text{ is divisible by } 3\}$ .
- $\{w \mid \text{every pair of adjacent } 0\text{'s in } w \text{ appears before any pair of adjacent } 1\text{'s}\}$ .
- $\{w \mid \text{every odd position of } w \text{ is a } 1\}$ .

**Problem 5.** Prove that the language  $\mathbf{L} = \{w \mid w \in \{0, 1\}^* \text{ is not a palindrome}\}$  is not regular.